

ITI 1121. Introduction to Computing II

Stack: applications

by

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Preamble

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Overview

Stack: applications

We look at several examples of the use of stacks, including evaluating arithmetic expressions, saving command history, and running Java programs.

General objective :

- ✚ This week you will be able to apply the stacks for algorithm design.

Preamble

Learning objectives

Learning objectives

- ✦ **Justify** the role of a stack in solving a computer problem.
- ✦ **Design** a computer program requiring the use of a stack.

Readings:

- ✦ Pages 159-176 of E. Koffman and P. Wolfgang.

Preamble

Plan

Plan

- 1 Preamble
- 2 Applications
- 3 Prologue

Applications

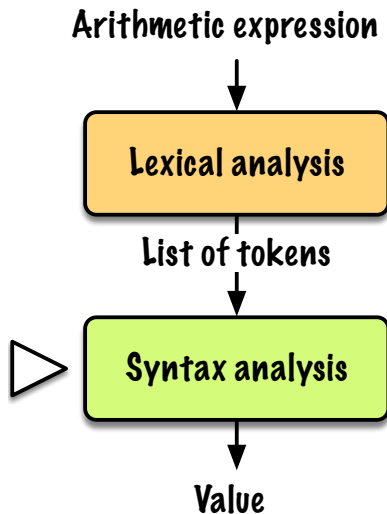
Applications

Evaluating an arithmetic expression

Application : Evaluating an arithmetic expression

Architecture of our application

- Clear separation of concerns: **lexical analysis** and **syntactic analysis**
- Lexical analysis** takes a **string of characters** as input and cuts it into chunks called **tokens**.
 - Input: $.1 \cdot + \cdot \cdot 2 \times 33 \cdot \cdot \cdot - 4 \cdot$
 - Output: $[1, +, 2, \times, 33, -, 4]$
- Our **syntax analysis** takes a sequence of **tokens** as input and returns the **value** of the expression.
 - Input: $[1, +, 2, \times, 33, -, 4]$
 - Output: 63



StringTokenizer

Java has a **lexical analyzer**!

```
StringTokenizer st;  
st = new StringTokenizer(" 1 + 2 * 33 - 4");  
  
while (st.hasMoreTokens()) {  
    System.out.println(st.nextToken());  
}
```

1
+
2
*
33
-
4

StreamTokenizer is more versatile!

Scan

Please take a few minutes to analyze this example. **What do you think?**

```
public static int scan(String expression) {  
    StringTokenizer st; String op; int l, r;  
  
    st = new StringTokenizer(expression);  
  
    l = Integer.parseInt(st.nextToken());  
  
    while (st.hasMoreTokens()) {  
        op = st.nextToken();  
        r = Integer.parseInt(st.nextToken());  
        l = eval(l, op, r);  
    }  
  
    return l;  
}
```

```
private static int eval(int l, String op, int r) {
    int result;
    switch (op) {
        case "+":
            result = l + r;
            break;
        case "-":
            result = l - r;
            break;
        case "/":
            result = l / r;
            break;
        case "*":
            result = l * r;
            break;
        default:
            System.exit(-1);
    }
    return result;
}
```

Exercises

- ❖ What does `scan("3 * 12 + 4")` return?
- ❖ What does `scan("3 + 12 * 4")` return?
- ❖ What do you think?

Discussion

- ❖ The **scan** algorithm evaluates operations from left to right, regardless of the **priority** of the operations. **Scan** doesn't process the **parentheses**.

There are two solutions:

- ❖ Use a new representation for the expressions
- ❖ Use a more complex algorithm

⇒ Both of these solutions require the use, implicit or explicit, of a **stack**!

Representations. There are three ways to represent an expression: $l \diamond r$, where \diamond is an operator.

infix: The infix notation corresponds to the usual notation, the operator is sandwiched between its operands: $l \diamond r$.

post-fixed: In post-fixed notation, the operands are placed in front of the operator, $l r \diamond$. This notation is also called *Reverse Polish Notation* or **RPN**, it is the notation used by some scientific calculators (such as HP-35 from Hewlett-Packard or Texas Instruments TI-89 using RPN Interface by Lars Frederiksen*) and the languages **PostScript** and **PDF**.

❖ $7 - (3 - 2) \rightarrow 7 3 2 - -$

❖ $(7 - 3) - 2 \rightarrow 7 3 - 2 -$

pre-fixed The third notation is to place the operator first followed by its operands, $\diamond l r$. The programming language **Lisp** uses a combination of parentheses and prefix notation, $(- 7 (* 3 2))$.

*www.calculator.org/rpn.html

From infix to postfix

- ❖ Successively transform, **one by one**, each subexpression **following the normal order of evaluation** of an infix expression.
- ❖ An infix subexpression $l \diamond r$ becomes $l r \diamond$, where l and r are themselves subexpressions and \diamond is an operator.

Evaluating a postfix expression (mentally)

Until the end of the expression is reached:

1. Read from **left to right** up to the **first operator**;
2. **Apply** the operator to the (2) operands preceding it.;
3. **Replace** the operator and its (2) operands by the result.

When the end of the expression is reached, we have the result.

Evaluating a postfix expression (mentally)

A few **exercises**:

✚ 9 3 / 10 2 3 * - +

✚ 9 2 4 * 5 - /

9 3 / 10 2 3 * - +

9 2 4 * 5 - /

Remarks

The **order of the operands is the same** for the two notations, postfix and infix, however the **places where the operators are inserted differ**.

❖ $2 + (3 * 4) \rightarrow 2 \ 3 \ 4 \ * \ +$

❖ $(2 + 3) * 4 \rightarrow 2 \ 3 \ + \ 4 \ *$

To evaluate an infix expression, the **operator priority** as well as the **parentheses** must be taken into account.

- ❖ In the case of postfix notation, **these concepts are represented within the notation**.

9 3 / 10 2 3 * - +

9 2 4 * 5 - /

Exercises

- Give the **content of the stack** for each iteration of the algorithm. :
 - 9 3 / 10 2 3 * - +
 - 9 2 4 * 5 - /
- Modify the algorithm so that it constructs an expression **infix** from an expression **postfix** given as input.

Applications

Discussion on the usefulness of abstract data types

Discussion

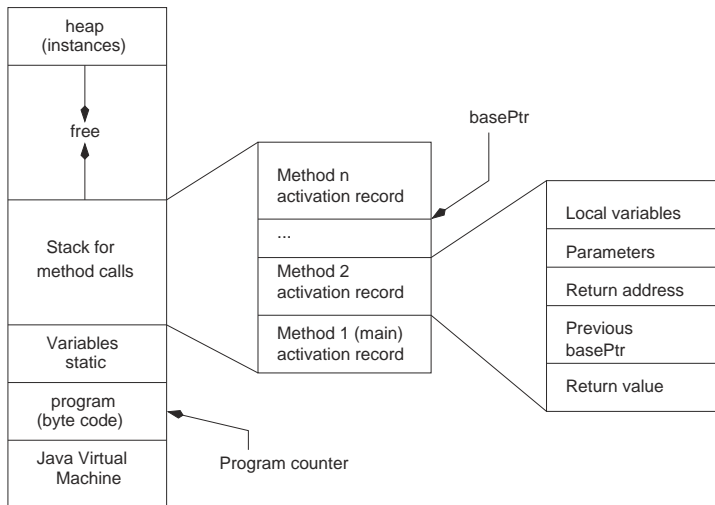
- ⊞ Now please answer the question asked earlier: “One of the proposed implementations uses an array, **why don't we just use an array** for algorithm design? **What are the advantages?** ”

Applications

Memory management

Application : Memory management
during program execution

Memory representation and program interpretation



When a method is called

The Java Virtual Machine (**JVM**) must:

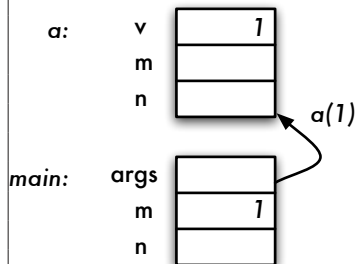
1. Create a new **activation frame** [working memory] (the return value, previous basePtr value and the return address have a fixed size, the size of local variables and parameters depends on the method);
2. **Save** the current value of basePtr, in the space “previous value of basePtr”, point basePtr to the base of the current block;
3. **Save** the value of locationCounter in the space designated by “return address”, make locationCounter point to the first instruction of the called method;
4. **Copy the actual parameter values in the region designated by “parameter”**;
5. Initialize the textbflocal variables;
6. **Start** execution at the instruction pointed to by locationCounter.

When the execution of a method ends

1. The method **saves the return value** at the location indicated by “return value”;
2. **Returns control** to the calling method, i.e., resets the locationCounter and basePtr values;
3. **Removes the current activation block**;
4. **Resumes** execution at the location designated by locationCounter.

```
public static int c(int v) {
    int n;
    n = v + 1;
    return n;
}
public static int b(int v) {
    int m,n;
    m = v + 1;
    n = c(m);
    return n;
}
public static int a(int v) {
    int m,n;
    m = v + 1;
    n = b(m);
    return n;
}
public static void main(String [] p) {
    int m = 1,n;
    n = a(m);
    System.out.println(n);
}
```

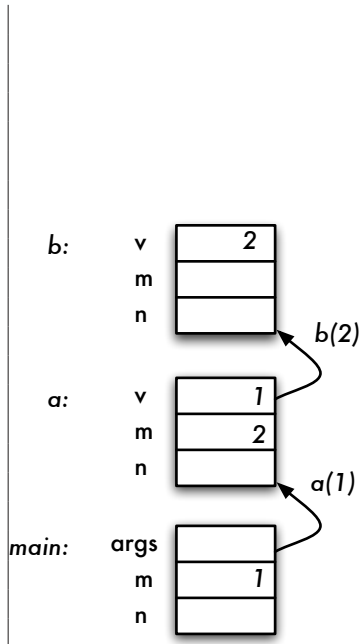
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public static int c(int v) {
    int n;
    n = v + 1;
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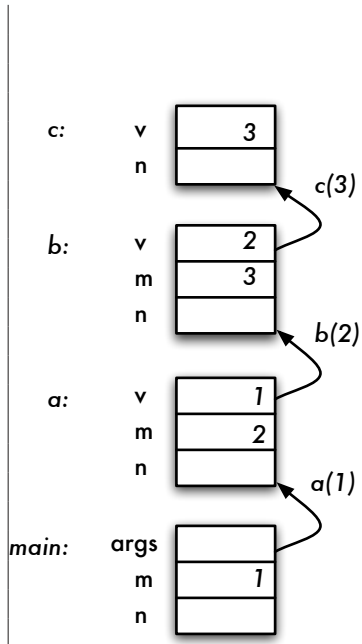
```



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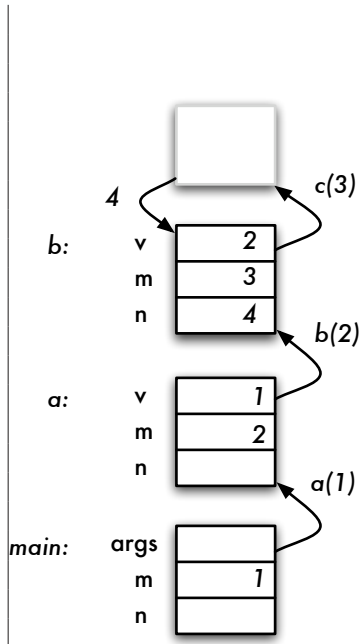
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public static void main(String[] p) {
    int m = 1,n;
    n = a(m);
    System.out.println(n);
}

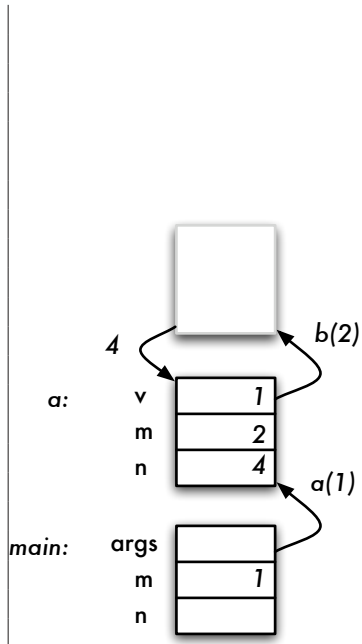
```



```

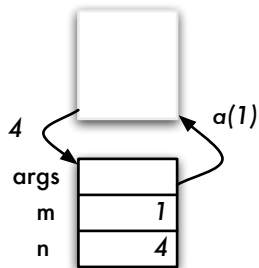
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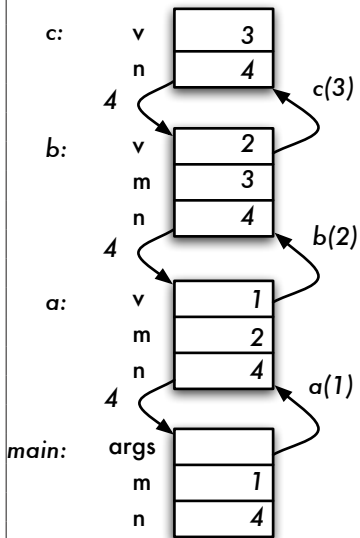
main:



```

public static int c(int v) {
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    return n;
}
public static int b(int v) {
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    m = v + 1;
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public static int a(int v) {
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    System.out.println(n);
}

```



Prologue

Summary

- ✚ A **stack** is used when one wishes to process the elements **in reverse order**.

Next module

- ▣ **Stack** : linked elements

References I



E. B. Koffman and Wolfgang P. A. T.

Data Structures: Abstraction and Design Using Java.

John Wiley & Sons, 3e edition, 2016.



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